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# Fibred Closed Braids with Disc-Band Fibre Surfaces. 

Marta Rampichini

Sunto. - Un risultato classico di Stallings fornisce una condizione necessaria e sufficiente per stabilire se una data superficie immersa senza autointersezioni in $S^{3}$ è una fibra per $S^{3}$ - $2 S$. In questo articolo si descrive come trovare una possibile fibra per un link presentato come treccia chiusa. Si descrive anche un algoritmo, implementato al calcolatore, che permette di trovare i principali ingredienti per verificare la condizione necessaria e sufficiente di Stallings, cioè una presentazione del gruppo fondamentale della superficie e del suo complementare in $S^{3}$, e una espressione esplicita dell'omomorfismo indotto in omotopia dalla mappa di push-off. L'articolo termina con una discussione di particolari proprietà della presentazione del gruppo $\pi_{1}\left(S^{3} \backslash S_{W}\right)$.

Summary. - A classical result by Stallings provides a necessary and sufficient condition to decide whether a given embedded surface $S$ is a fibre in $S^{3} \backslash \partial S$. In this paper it is described how to find a candidate fibre surface for a a link presented as a closed braid. Also it is described an implemented algorithm to find the main ingredients of the necessary and sufficient condition of Stallings, namely presentations of the fundamental groups of the surface and of its complement in $S^{3}$, and an explicit expression of the homomorphism induced in homotopy by the push-off map. The paper ends with a discussion of the particular properties of the presentation of $\pi_{1}\left(S^{3} \backslash S_{W}\right)$.

## 1. - Introduction.

Fibred three-manifolds are very nice topological objects: they can be described as a cartesian product of a surface cross an interval, $S \times[0,1]$, modulo a gluing map $h$ which identifies the two copies of the surface, $h: S_{0} \rightarrow S_{1}$. When a surface $S$ is embedded in $S^{3}$ so that the complement of its boundary is a three-manifold fibred over $S^{1}$, then $S$ is said to be a fibre surface, and its boundary is called a fibred link. The unknot is fibred, since its complement in $S^{3}$ can be filled by a continuous family of discs, all with the given unknot as boundary. Famous examples of fibred links are the algebraic links [13] and the closure of homogeneous braids [24]. Harer gave instructions for constructing all possible fibred links starting from the unknot, by modifying their fibre surfaces [10].

Montesinos and Morton constructed all fibred links by using branched coverings on closed braids [14].

Gabai gave a necessary and sufficient topological condition for a link to be fibred by looking at a candidate fibre surface [7]. The question of detecting whether a given link is fibred is related to the problem of finding a candidate fibre surface. An algorithm for finding minimal genus fibred surfaces, based on normal surfaces with respect to a triangulation of the three manifold, is given in $[25,16]$.

In this paper, I am particularly interested in finding fibre surfaces for links presented as closed braids, using the explicit expression of the braid for finding a fibre surface, looking for direct connections between algebraic properties of the braid and topological properties of the surface.

Some necessary conditions for a Seifert surface of a link to be a fibre are summarized in the following:

Theorem 1 (cf. [23, 6] Prop. 4.1). - If S is a fibre surface then it is connected and it is a Seifert surface of minimal genus for its boundary. If $L$ is a fibred link, then all its minimal genus Seifert surfaces are fibre surfaces and they are all isotopic.

There are partial results about the problem of finding minimal genus Seifert surfaces for particular links, see [8, 9, 26], but for general links the problem is still open.

A classical result by Stallings [24] gives a necessary and sufficient algebraic condition for a given surface to be a fibre. It involves the fundamental groups of the surface and of its complement in $S^{3}$. It consists in deciding whether the particular homomorphism induced in homotopy by the push-off map is an isomorphism. However, if we want to use this result, we first need to find a candidate fibre surface.

Another way to attack the problem is that of representing links as closed braids, which allows to deal with them as algebraic objects: this is not completely equivalent to the previous setting, since any link can be represented as a closed braid in infinitely many different ways. But if we fix a braid axis for the link, then we will be dealing with a single conjugacy class in a specified braid group [15].

In order to make the paper self-contained, I have recalled known results where useful. In order to make my algorithm more clear, I have considered it along with an example, throughout the paper.

The paper is organized as follows.
In Section 2 some interesting properties of closed braids are collected, and some Seifert surfaces embedded in a particular way are shown. I
will define a fibred closed braid and in Proposition 5 I will prove how to find a candidate fibre surface for a closed braid.

In Section 3 I explain an algorithm to find the ingredients of Stallings theorem, namely: presentations for the fundamental groups of the surface and of its complement in $S^{3}$, and an explicit expression for the homomorphism. Also I prove that the homomorphism is always injective.

In Section 4 I explain some properties of the presentation of $\pi_{1}\left(S^{3} \backslash S\right)$, in order to attack the problem (which in general is not solvable) of determining whether the homomorphism is an isomorphism.

In the Appendix I explain an algorithm for computing the Alexander polynomial for a braid given as a word in disc-band generators.

Throughout the paper, all topological objects under consideration (links, surfaces, braids) are oriented.

The algorithms presented in this paper have been implemented using GAP, a programming language for Group Theory $\left({ }^{1}\right)$, with the help of P. Boldi $\left({ }^{2}\right)$ and S. Vigna $\left({ }^{(2)}\right.$.

## 2. - Fibred closed braids.

The band presentation of the braid group given by Birman, Ko and Lee in [2] provides a translation of many combinatorial and algebraic properties of braid words and braids into topological properties of disc-band surfaces, which are Seifert surfaces for the closed braids, embedded in a particular way.

In this section I will recall the band presentation of $B_{n}$, the definition of disc-band surface, and some useful properties.

The band generators are a generalisation of the classical Artin generators: for $B_{n}$ there are $\binom{n}{2}$ generators $a_{j, i}$ with $1 \leqslant j<i \leqslant n$, where $a_{j, i}$ is the braid in which the $i^{\text {th }}$ strand crosses over the $j^{\text {th }}$ strand, both going over the other intervening strands. The set of band generators includes the set of classical Artin generators, in fact $\sigma_{i}=a_{i, i+1}$. To simplify the notation, I will write $(j, i)$ for $a_{j, i}$ and $\overline{(j, i)}$ for $a_{j, i}^{-1}$.

There are two types of relations (see Fig. 1):

1. $(j, i)(k, h)=(k, h)(j, i)$ whenever $(i-h)(i-k)(j-h)(j-k)>0$, the condition meaning that the two pairs of indices are not interleaved, and
2. $(j, i)(k, j)=(k, j)(k, i)=(k, i)(j, i)$ whenever $n \geqslant i>j>k \geqslant 1$.
${ }^{(1)}$ [GAP 99] The GAP Group, GAP - Groups, Algorithms, and Programming, Version 4.1; Aachen, St Andrews, 1999. The package is available at http://www-gap.dcs.stand.ac.uk/ gap/gap.html.
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Fig. 1. - Relations of band generators.

To each word in band generators we can associate a Seifert surface for the closed braid embedded in a particularly nice way.

Definition 2. - Given a word $W$ expressed in the band generators of $B_{n}$, we associate to it the disc-band surface $S_{W}$, constructed in the following way:

- take $n$ discs $D_{k}$, each pierced once by the braid axis, all in the same direction;
- for any letter $a_{i, j}^{\varepsilon}$ of $W$, connect the discs $D_{i}$ and $D_{j}$ by an half-twisted band, the twist depending on the sign of the letter.

If $\beta$ is the braid represented by $W$, the boundary of $S_{W}$ is the closure $\widehat{\beta}$, and the sign of each letter corresponds to the sign of the associated crossing.

See Fig. 2 for an example. I will go on with this example throughout the paper.

Disc-band surfaces are also called braided surfaces in [19, 20, 21, 22]. Following Rudolph [19] let me represent each half-twisted band as a half curl, so that we always see the positive side of the surface, as in Fig. 3.

There are infinitely many disc-band surfaces associated to a given closed braid: in fact there are closed braids with more than three strands that can be represented by infinitely many different conjugacy classes of braids [4]; moreover any braid can always be represented by infinitely many different words in the band generators. Nevertheless, many algebraic


Fig. 2. - The disc-band surface $S_{W}$ for the braid word $W=(5,6)(1,4)$ $\overline{(2,5)(3,4)}(1,6)(1,3) \overline{(2,5)}(3,6)$.
and combinatorial equivalences of braids and braid words can be easily translated into isotopies of these surfaces.

The following algebraic changes in a braid word $W$ correspond to isotopies of the associated disc-band surface $S_{W}$ in the complement of the braid axis:

1. the disc-band relations; an example of the corresponding isotopy is shown in Fig. 4;
2. conjugations by $\delta=(n-1, n)(n-2, n-1) \ldots(1,2)$, which have the effect of shifting indices by $1(\bmod n)$ on each band generator: $\delta^{-1}(j, i) \delta=$ $(j+1, i+1)$; the corresponding isotopy is a cycling of the discs $D_{k}$ of $S_{W}$ along the axis;


Fig. 3. - How to see the half-twisted bands as half curls.


Fig. 4. - How to see the relation $(j, i)(k, j)=(k, j)(k, i)$ as isotopy of disc-band surfaces.
3. conjugations by initial or final letters of $W$, which correspond to cycling the letters of $W$; the corresponding isotopy is a cicling of the bands around the dises of $S_{W}$.

I call 2 and 3 easy conjugations.
It is still open the question whether all isotopies among two disc-band surfaces in the complement of the axis can be expressed by a finite chain of moves 1, 2 and 3 above, cf Rudolph [22], p. 263.

We can also read from a word $W$ an important topological property of the surface $S_{W}$. The Euler characteristic of $S_{W}$ is related to the number $n$ of strands of the braid and the length $l(W)$ of the word in the band generators:

$$
\chi\left(S_{W}\right)=n-l(W)
$$

this can be seen for instance by retracting $S_{W}$ on a graph $\Gamma_{W}$ as shown in figures 7 and 8.

Notice that all relations and easy conjugations are length preserving. Instead, the free reduction of words is not length preserving, and in fact does not correspond to an isotopy of disc-band surfaces.

In what follows, any link $L$ will be seen as a closed braid $\widehat{\beta}$, with both the axis $A$ and the number $n=l k(L, A)$ of strands fixed. This will allow me to study $L$ as an algebraic object, the conjugacy class [ $\beta$ ] in $B_{n}$. This will not compromise my study, thanks to the following:

Theorem 3 (Rudolph, Proposition 1 of [20]). - For any fixed braid axis $A$ (i.e. for any fixed embedded unknot, with the fibration by discs of $S^{3} \backslash A$ fixed), any given embedded oriented connected fibre surface $S$ can be isotoped in $S^{3}$ so that it lies as a disc-band surface relative to the chosen braid axis $A$.

In what follows the term «closed braid» will mean both the isotopy class of a closed braid in $S^{3} \backslash A$, or the corresponding conjugacy class in $B_{n}$.

Definition 4. - A closed braid $\widehat{\beta}$ is called a fibred closed braid if it has a disc-band surface which is a fibre.

Proposition 5. - A closed braid $\widehat{\beta}$ is fibred if and only if for each word $W$ of minimal length in the conjugacy class of $\beta$ the surface $S_{W}$ is a fibre surface.

Proof. - If $\widehat{\beta}$ is fibred, then all its fibres are isotopic and have minimal genus, and by definition there exists at least one which is a disc-band surface. Since the genus of a disc-band surface is minimal (among disc-band surfaces) when the length of the word is minimal in the conjugacy class, there is at least one of these minimal-length words such that $S_{W}$ is a fibre. But for all such words the corresponding surfaces are isotopic, since they have the same genus and isotopic boundary; therefore they all are fibre surfaces.

Notice that any fibred link $L$ can be seen as a fibred closed braid $\widehat{\beta}$, with $\beta \in B_{n}$, but usually it is not known for which $n$ this is possible. The theorem of Bennequin studied in [3] guarantees that all braids with at most three strands have minimal genus disc-band Seifert surfaces, but this is no longer true in $B_{n}$ with $n \geqslant 4$ : Ko and Lee in [11] give an example by Morton of a 4-braid whose closure is the unknot (therefore it is a fibred link), which has not a disc-band surface of minimal genus, because the minimal length of a word in its conjugacy class is 7 , while it should be 3 . Anyway in $B_{4}$ there are conjugacy classes representing the unknot and with minimal length 3.

So, if we are given a braid and we want to find a candidate fibre surface we have to look for a minimal-length word in its conjugacy class, whose disc-band surface is connected.

The general problem of finding a minimal-length word in a given conjugacy class of $B_{n}$ is not solved, but for fibred closed braids we can use the Alexander polynomial of the closed braid to know what the minimal length should be.

Proposition 6. - The following are necessary conditions for a closed braid $\widehat{\beta}$ to be fibred:

1. the minimal length $l$ for a word in its conjugacy class is such that

$$
\begin{aligned}
& l \geqslant \operatorname{deg} \Delta_{\widehat{\beta}}(t)+n-1 \text {, if } \widehat{\beta} \text { has one component, and } \\
& l \geqslant \operatorname{deg} \Delta_{\widehat{\beta}}(t)+n \text {, if } \widehat{\beta} \text { has more than one component }
\end{aligned}
$$

2. in the Alexander polynomial $\Delta_{\widehat{\beta}}(t)$ the coefficients of the extreme powers of $t$ are $\pm 1$.

Proof. - This is a corollary of the following:
Theorem 7 (cf. [17, 18]). - A necessary condition for a link $L$ to be fibred is that the Alexander polynomial $\Delta_{L}(t)$ satisfies the following conditions:

1. if $S$ is a fibre, then

$$
\operatorname{deg} \Delta_{L}(t) \leqslant 1-\chi(S) \text {, if } L \text { has one component; }
$$

$\operatorname{deg} \Delta_{L}(t) \leqslant-\chi(S)$, if $L$ has more than one component;
2. in the Alexander polynomial $\Delta_{L}(t)$ the coefficients of the extreme powers of $t$ are $\pm 1$.

In the Appendix I explain an algorithm for computing the Alexander polynomial of a braid expressed by a word in band generators.

Once we have computed the Alexander polynomial, we can easily check the required property on the coefficients of extreme powers of $t$. If this is satisfied, we can look for a word with prescribed length and with the same exponent sum as $\beta$; then we can check conjugacy e.g. by the algorithm of [2]. The set of words of $B_{n}$ with prescribed length and exponent sum is finite, so either we end with a negative answer, or we find a word $W$ giving a candidate fibre surface $S_{W}$, to which we will apply the Stallings theorem.

## 3. - The algorithm.

The aim of this section is to explain how to find in our settings the main ingredients of Stallings' necessary and sufficient condition for a given surface to be a fibre. We first need the following:

Definition 8. - Given a connected oriented surface $S$ embedded in $S^{3}$, the push-off map $i^{+}: S \rightarrow S^{3} \backslash S$ is the continuous map (defined up to isotopy in $\left.S^{3} \backslash \partial S\right)$ which sends $S$ in a homeomorphic copy $S^{+}$, which lies in $S^{3} \backslash S$, obtained from $S$ by pushing each point off from $S$ along the positive normal direction.

Theorem 9 (Stallings, [24]). - A given connected surface $S$ embedded in $S^{3}$ is a fibre in $S^{3} \backslash \partial S$ if and only if the homomorphism

$$
i_{*}^{+}: \pi_{1}(S) \rightarrow \pi_{1}\left(S^{3} \backslash S\right)
$$

induced in homotopy by the push-off map is an isomorphism.

- The fundamental group $\pi_{1}\left(S_{W}\right)$ is a free group on $1-\chi\left(S_{W}\right)$ generators, but we need an explicit expression of such generators in topological terms to


Fig. 5. - The surface $S_{W}$ of Fig. 2, flattened, with the vertices of the graph $\Gamma_{W}$ embedded as intersections of $S_{W}$ with the braid axis. The edge $e_{7}$ is shown as an example. Each generator $\beta_{j}$ of $\pi_{1}\left(S^{3} \backslash S_{W}\right)$ is a loop from your eye to the arrow and back to your eye.
know how the homomorphism $i_{*}^{+}$will act on them. We can read them explicitly from $W$ : look at Fig. 5, where the surface of Fig. 2 has been flattened as explained in Fig. 3.

To find an explicit expression of the generators, we can consider a graph $\Gamma_{W}$ which is a strong deformation retract of $S_{W}$, look at a spanning tree $T_{W}$ of $\Gamma_{W}$ and then read a loop $\alpha_{k}$ for each edge in $\Gamma_{W} \backslash T_{W}$. All this information can be read directly from $W$. Choose a base point on $D_{1}$ in a small neighbourhood of the point where the braid axis intersects the disc. Suppose $W=\prod_{j=1}^{l(W)} w_{j}^{\varepsilon_{j}}$,


Fig. 6. - The graph $\Gamma_{W}$ (a) and a spanning tree $T_{W}$ for it (b).
where each $w_{j}$ is a band generator and $\varepsilon_{j}= \pm 1$. The graph $\Gamma_{W}$ has one vertex $v_{j}$ on each disc $D_{j}$ of $S_{W}$, that is one vertex for each index $j \in\{1,2, \ldots n\}$ (they are marked by «big dots» in Fig. 5). The edges of $\Gamma_{W}$ are one for each band of $S_{W}$, that is one edge $e_{i}$ for each letter $w_{i}$. If $w_{i}=(h, k)$, then $e_{i}$ is a simple arc which joins $v_{h}$ to $v_{k}$ going along $D_{h}$, then along the $i^{\text {th }}$ band then along $D_{k}$ (in Fig. 5, the edge $e_{7}$ is shown as an example). We can see a schematic picture of $\Gamma_{W}$ in Fig. 6 (a) and a spanning tree in Fig. 6 (b). Our GAP procedures Build$\operatorname{Graph}(W, n)$ and $\operatorname{SpTree}(G)$ give automatically the graph and the spanning tree from the braid word $W$.

It is convenient for the implementation of the algorithm to write the $1-$ $\chi\left(S_{W}\right)$ generators $\alpha_{h}$ 's as words in the letters $e_{i}$ 's, with the convention that each edge is directed from the smallest to the greatest of its two vertices. This is done by our GAP procedure ReadGenerators $(G, T)$, which for our example gives

$$
\begin{aligned}
& \alpha_{1}=e_{6} e_{4} e_{2}^{-1}, \\
& \alpha_{2}=e_{5} e_{1}^{-1} e_{3}^{-1} e_{7} e_{1} e_{5}^{-1}, \\
& \alpha_{3}=e_{6} e_{8} e_{5}^{-1}
\end{aligned}
$$

- We can apply the well known Wirtinger algorithm to find a presentation of the fundamental group $\pi_{1}\left(S^{3} \backslash S_{W}\right)$. Look at Fig. 5: choose one generator $\beta_{i}$ for each band of $S_{W}$, i.e. for each letter $w_{i}$ of $W$; each generator is a loop surrounding exactly one band, as indicated schematically by the arrow in Fig. 5. To read the relations, the surface can be retracted as in Fig. 7 and 8.

In order to write the relations, for each generator $\beta_{k}$ I define two conjugates, $\gamma_{k}$ and $\delta_{k}$, which can be seen in the corner of Fig. 8: $\gamma_{k}$ (respectively $\delta_{k}$ ) is the loop on the same strand as $\beta_{k}$, but as near as possible to the initial (respectively final) vertex of the simple arc $e_{k}$. Their expression as conjugates of
(a)

(b)


Fig. 7. - (a) The disc $D_{1}$ of Fig. 5 and (b) its retract.
the $\beta_{k}$ 's can be obtained from $W$ as follows: suppose $w_{k}=(i, j)$ and let $w_{r}=$ ( $p, q$ ) be another generic letter. Then

$$
\gamma_{k}=P_{k} \beta_{k} P_{k}^{-1} \text { and } \delta_{k}=Q_{k} \beta_{k} Q_{k}^{-1}
$$



Fig. 8. - The graph $\Gamma_{W}$ as deformation retract of the surface $S_{W}$, with arrows indicating the generators of $\pi_{1}\left(S^{3} \backslash S_{W}\right)$. In a corner, a zoom of a neighbourhood of a vertex.
where

$$
P_{k}=\prod_{r=1}^{k-1} \beta_{r}^{s_{r}} \quad \text { and } \quad Q_{k}=\prod_{r=1}^{k-1} \beta_{r}^{t_{r}},
$$

and

$$
s_{r}=\left\{\begin{array}{ll}
1 & \text { if } p \leqslant i<q \\
0 & \text { otherwise },
\end{array} \quad t_{r}= \begin{cases}1 & \text { if } p \leqslant j<q \\
0 & \text { otherwise }\end{cases}\right.
$$

For instance, from Fig. 8 it is easy to read $\gamma_{4}=\beta_{2} \beta_{3} \beta_{4} \beta_{3}^{-1} \beta_{2}^{-1}$.
The conjugates $\gamma_{k}$ and $\delta_{k}$ are useful in order to read relations of the group $\pi_{1}\left(S^{3} \backslash S_{W}\right)$; look at a neighbourhood of each vertex (as at the corner of Fig. 8): we get

$$
r_{m}=\prod_{k=1}^{l} \gamma_{k}^{x_{k}} \delta_{k}^{y_{k}},
$$

where

$$
x_{k}=\left\{\begin{array}{cl}
-1 & \text { if } i=m \\
0 & \text { otherwise }
\end{array} \quad y_{k}=\left\{\begin{array}{cl}
1 & \text { if } j=m \\
0 & \text { otherwise }
\end{array}\right.\right.
$$

For instance, from Fig. 8 it is easy to read $r_{3}=\gamma_{4}^{-1} \delta_{6} \gamma_{8}^{-1}$.
Our GAP procedure CompleteHomotopyRelations $(F, W, n)$ (here $F$ is the free group on $l(W)$ generators) gives

$$
\begin{aligned}
& r_{1}=\beta_{6}^{-1} \beta_{5}^{-1} \beta_{2}^{-1} \\
& r_{2}=\beta_{2} \beta_{5} \beta_{6} \beta_{7}^{-1} \beta_{6}^{-1} \beta_{5}^{-1} \beta_{3}^{-1} \beta_{2}^{-1} \\
& r_{3}=\beta_{2} \beta_{3} \beta_{5} \beta_{6} \beta_{7} \beta_{8}^{-1} \beta_{7}^{-1} \beta_{5}^{-1} \beta_{4}^{-1} \beta_{3}^{-1} \beta_{2}^{-1}, \\
& r_{4}=\beta_{2} \beta_{3} \beta_{4} \beta_{3}^{-1} \\
& r_{5}=\beta_{3} \beta_{5} \beta_{7} \beta_{5}^{-1} \beta_{1}^{-1} \\
& r_{6}=\beta_{1} \beta_{5} \beta_{8}
\end{aligned}
$$

There exists a GAP procedure which simplifies presentations (using Tietze's moves); in our example, it immediately gives a free group on three generators. We get $\pi_{1}\left(S^{3} \backslash S_{W}\right)=\left\langle\beta_{2}, \beta_{3}, \beta_{5} \mid-\right\rangle$, with

$$
\begin{aligned}
& \beta_{1}=\beta_{3} \beta_{5} \beta_{2} \beta_{3}^{-1} \beta_{2}^{-1} \beta_{5}^{-1} \\
& \beta_{4}=\beta_{3}^{-1} \beta_{2}^{-1} \beta_{3} \\
& \beta_{6}=\beta_{5}^{-1} \beta_{2}^{-1} \\
& \beta_{7}=\beta_{2} \beta_{3}^{-1} \beta_{2}^{-1} \\
& \beta_{8}=\beta_{2} \beta_{3} \beta_{2}^{-1} \beta_{5}^{-1} \beta_{3}^{-1}
\end{aligned}
$$

- The homomorphism $i_{*}^{+}$. For each arc $e_{k}$ one can see in Fig. 5 how to read its image $i^{+}\left(e_{k}\right)$ as a product of loops $\beta_{j}$ 's: it is sufficient to read one $\beta_{j}$ each time $e_{k}$ passes under the corresponding band in the same direction as $\beta_{j}$, and the inverse when it passes in the opposite direction. All this information can be read directly from $W$ as follows:

$$
i^{+}\left(e_{k}\right)=P_{k} \beta_{k}^{\varepsilon} Q_{k}^{-1}
$$

where $\varepsilon=\frac{\varepsilon_{k}+1}{2}$ and $P_{k}, Q_{k}$ have been defined above. Our GAP procedure IPlusEks $(W)$ computes the list of images of the $e_{k}$ 's. In our example we get

$$
\begin{aligned}
& i^{+}\left(e_{1}\right)=\beta_{1} \\
& i^{+}\left(e_{2}\right)=\beta_{2} \\
& i^{+}\left(e_{3}\right)=\beta_{2} \beta_{1}^{-1} \\
& i^{+}\left(e_{4}\right)=\beta_{2} \beta_{3} \beta_{3}^{-1}=\beta_{2} \\
& i^{+}\left(e_{5}\right)=\beta_{2} \beta_{5} \\
& i^{+}\left(e_{6}\right)=\beta_{2} \beta_{5} \beta_{6} \beta_{5}^{-1} \beta_{4}^{-1} \beta_{3}^{-1} \beta_{2}^{-1} \\
& i^{+}\left(e_{7}\right)=\beta_{2} \beta_{3} \beta_{5} \beta_{6} \beta_{5}^{-1} \beta_{1}^{-1} \\
& i^{+}\left(e_{8}\right)=\beta_{2} \beta_{3} \beta_{4} \beta_{5} \beta_{7} \beta_{8} .
\end{aligned}
$$

Now we can use the expression of the generators $\alpha_{j}$ as products of $e_{k}$ 's and substitute their images $i^{+}\left(e_{k}\right)$ to get the images $i_{*}^{+}\left(\alpha_{j}\right)$. Our GAP procedure $\operatorname{IPlus}(W, n, H)$ gives the list of images $i_{*}^{+}\left(\alpha_{j}\right)$, using the word $W$, the braid index $n$ and the fundamental group $H=\pi_{1}\left(S^{3} \backslash S_{W}\right)$ as found by our GAP procedure HomotopyGroup $(W, n)$. In our example, we get:

$$
\begin{aligned}
& i_{*}^{+}\left(\alpha_{1}\right)=\beta_{2} \beta_{5} \beta_{6} \beta_{5}^{-1} \beta_{4}^{-1} \beta_{3}^{-1} \beta_{2}^{-1} \\
& i_{*}^{+}\left(\alpha_{2}\right)=\beta_{2} \beta_{5} \beta_{3} \beta_{5} \beta_{6} \beta_{5}^{-2} \beta_{2}^{-1} \beta_{5}^{-2} \beta_{2}^{-1} \\
& i_{*}^{+}\left(\alpha_{3}\right)=\beta_{2} \beta_{5} \beta_{6} \beta_{7} \beta_{8} \beta_{5}^{-1} \beta_{2}^{-1} .
\end{aligned}
$$

In this case, by substituting the expressions found above we get

$$
\begin{aligned}
& i_{*}^{+}\left(\alpha_{1}\right)=\beta_{5}^{-1} \beta_{3}^{-1} \\
& i_{*}^{+}\left(\alpha_{2}\right)=\beta_{2} \beta_{5} \beta_{3} \beta_{2}^{-1} \beta_{5}^{-2} \beta_{2}^{-1} \beta_{5}^{-2} \beta_{2}^{-1} \\
& i_{*}^{+}\left(\alpha_{3}\right)=\beta_{5}^{-1} \beta_{3}^{-1} \beta_{5}^{-1} \beta_{2}^{-1},
\end{aligned}
$$

from which we can easily get

$$
\begin{aligned}
& \beta_{1}=i_{*}^{+}\left(\alpha_{3}^{-1} \alpha_{1} \alpha_{3}^{-1} \alpha_{1} \alpha_{3} \alpha_{2} \alpha_{3}^{-1} \alpha_{1}\right) \\
& \beta_{2}=i_{*}^{+}\left(\alpha_{1}^{-1} \alpha_{3}^{-1} \alpha_{1} \alpha_{3} \alpha_{2} \alpha_{3}^{-1} \alpha_{1}\right) \\
& \beta_{3}=i_{*}^{+}\left(\alpha_{1}^{-1} \alpha_{3} \alpha_{2}^{-1} \alpha_{3}^{-1} \alpha_{1}^{-1} \alpha_{3}\right) .
\end{aligned}
$$

To summarise, the algorithm presented in this section allows to find explicitly all the ingredients to apply Stallings' theorem, which are:

1) an expression of the generators of the free group $\pi_{1}\left(S_{W}\right)$ in terms of loops on $S_{W}$;
2) a presentation of the group $\pi_{1}\left(S^{3} \backslash S_{W}\right)$;
3) an explicit expression of the homomorphism $i_{*}^{+}$.

One more step can be done:
Proposition 10. - If $S_{W}$ is a connected surface of minimal genus for $\partial S_{W}$, then $i_{*}^{+}: \pi_{1}\left(S_{W}\right) \rightarrow \pi_{1}\left(S^{3} \backslash S_{W}\right)$ is injective.

Proof. - Since $S_{W}$ is of minimal genus, it is incompressible; this means what follows: for any simple closed curve $\gamma \subset S_{W}$, boundary of a disc $D \subset S^{3}$, with $D \cap S_{W}=\gamma$, there is a disc $D^{\prime} \subset S_{W}$ with $\partial D^{\prime}=\gamma$. Since $i_{*}^{+}([\gamma])=1$ in $\pi_{1}\left(S^{3} \backslash S_{W}\right)$ if and only if $i^{+}(\gamma)$ bounds a disc $D$ embedded in $S^{3} \backslash S_{W}$, one can attach to this disc the annulus $\boldsymbol{n} \times \gamma$ (where $\boldsymbol{n}$ is the normal positive vector used to define $i^{+}$): in this way one gets a disc in $S^{3}$ bounded by $\gamma$, therefore $\gamma$ is nullhomotopic in $S_{W}$.

## 4. - Combinatorial properties of the presentation of $\pi_{1}\left(S^{3} \backslash S_{W}\right)$.

To decide whether a braid is fibred or not, once we have found all ingredients of Stallings' theorem, we need to decide whether the homomorphism $i_{*}^{+}$is surjective or not. This is an open problem of combinatorial group theory, which is known to be solvable in some large families of groups: the free groups for instance, and the groups satisfying the so called small cancellation conditions. A good reference is [12].

To give an idea about the small cancellation conditions we need some terminology of combinatorial group theory (cf. [12]). Suppose we are given a presentation $\langle X \mid R\rangle$ for a group $G$, where $X$ is a set of generators and $R$ a set of relations. Suppose $w=x_{1} x_{2} \ldots x_{k}$ is a word in the generators (and their inverses): say that $w$ is cyclically reduced if $x_{k} \neq x_{1}^{-1}$. Say that $R$ is symmetrized if all elements of $R$ are cyclically reduced and, for each $r \in R$, all cyclically reduced conjugates of $r$ and of $r^{-1}$ belong to $R$. Now suppose that there are two elements $r_{1}, r_{2}$ of $R$ such that there are words $b, c_{1}, c_{2}$ such that $r_{1}=b c_{1}, r_{2}=$
$b c_{2}$. Then $b$ is called a piece relative to $R$. The small cancellation conditions assert that pieces are relatively small. The most usual conditions are the following:

- Condition $C^{\prime}(\lambda)$ for all $r \in R$, if $r=b c$ and $b$ is a piece, then $|b|<$ $\lambda|r|$, with $\lambda \in \boldsymbol{R}, \lambda>0$.
- Condition $C(p)$ : for all $r \in R, r$ is the product of at least $p$ pieces.

REmARK. - In both groups with free presentation and groups with presentations satisfying these small cancellation conditions for appropriate $\lambda, p$, the word problem is solvable.

In what follows I explain some combinatorial properties of the presentation of $\pi_{1}\left(S^{3} \backslash S\right)$ found in the preceding section.

Proposition 11. - For a given word $W \in B_{n}$, of minimal length $l=l(W)$ in its conjugacy class, the presentation for $\pi_{1}\left(S^{3} \backslash S_{W}\right)$ has $l$ generators and $n$ relations; each relation is the product of some conjugates $\gamma_{j}, \delta_{j}$ of the generators, with the following properties:

1) in each $\gamma_{j}=P_{j} \beta_{j} P_{j}^{-1}$, the generator $\beta_{j}$ occurs only once; and a generator $\beta_{k}$ can occur, at most once in $P_{j}$, only if $k<j$; similarly for $\delta_{j}=Q_{j} \beta_{j} Q_{j}{ }^{-1}$;
2) if $w_{j}=(p, q)$, then the length of $\gamma_{j}$ equals the number of letters $w_{h}$ with $h<j$ such that $w_{h}=(r, s), r \leqslant p<s$; similarly, the length of $\delta_{j}$ equals the number of letters $w_{h}$ with $h<j$ such that $w_{h}=(r, s), r \leqslant q<s$;
3) considering the set of $n$ relations, each $\gamma_{j}$ and each $\delta_{j}$ occur in one relation, and they do not occur in the same relation;
4) the number of $\gamma_{j}$ 's and $\delta_{j}$ 's occurring in $r_{k}$ equals the number of letters $w_{j}$ in which the index $k$ occurs.

Summarising we get for all $j=1,2, \ldots l$ and all $k=1,2, \ldots n$

$$
\begin{aligned}
& \left|\gamma_{j}\right| \leqslant 2 j+1 \\
& \left|\delta_{j}\right| \leqslant 2 j+1 \\
& \left|r_{k}\right|<l(2 l+1) .
\end{aligned}
$$

All these properties are clear from the explicit construction given in the preceding section.

Moreover, suppose there is one index $1 \leqslant i \leqslant n$ which only occurs in one letter of $W$. Then we can conjugate $W$ by its initial letter until the letter with the index $i$ becomes the first. Thus $r_{i}=\beta_{1}$, which means $\beta_{1}=1$, so it can be cancelled. This corresponds to the result of a Markov move on the surface.

Therefore we can always suppose that each index $1 \leqslant i \leqslant n$ occurs at least twice in $W$; hence for all relations we get

$$
\left|r_{i}\right| \geqslant 2
$$

Concerning the number of generators, we can say the following:
Proposition 12. - If $W$ is a word of minimal length in the conjugacy class of $\beta$

1) if $|W|=n-1$ and $\widehat{\beta}$ has one component, then $\widehat{\beta}$ is a fibred braid;
2) if $|W|=n-1$ and $\widehat{\beta}$ has more than one component, or if $|W|<n-$ 1 , then $\widehat{\beta}$ is neither a fibred braid nor a fibred link.

Proof. - In the first case $\chi\left(S_{W}\right)=1$, therefore $S_{W}$ is a disc, which is trivially a fibre surface. In the second case $S_{W}$ is a minimal genus Seifert surface which is not connected (for a surface $S=S_{1} \amalg S_{2}, \chi(S)=\chi\left(S_{1}\right)+\chi\left(S_{2}\right)$ ), therefore the boundary link cannot be fibred.

Corollary 13. - In the presentation of $\pi_{1}\left(S^{3} \backslash S_{W}\right)$ obtained as above, the number of generators is greater than or equal to the number of relations.

Proof. - We can always suppose $|W| \geqslant n$.
An interesting example. For the braid

$$
W=(1,4)(2,6)(1,3)(2,4)(3,5)(1,4)(2,6)(1,3)(2,4)(3,5)
$$

the Alexander polynomial

$$
\Delta_{\widehat{\beta}}(t)=-1+t^{2}-t^{4}
$$

satisfies the necessary conditions requested by Theorem 7. After some Tietze transformations, we get the presentation

$$
\pi_{1}\left(S^{3} \backslash S_{W}\right)=\left\langle\beta_{1}, \beta_{2}, \ldots, \beta_{6} \mid r\right\rangle
$$

where the relation is

$$
r=\beta_{2} \beta_{6}^{-1} \beta_{5}^{-1} \beta_{4}^{-1} \beta_{2}^{-1} \beta_{1}^{-1} \beta_{2} \beta_{5} \beta_{2}^{-1} \beta_{1} \beta_{2} \beta_{3} \beta_{4} \beta_{6} \beta_{2}^{-1} \beta_{6}^{-1} \beta_{3}^{-1} \beta_{1}^{-1}
$$

with length $|r|=18$, and pieces of maximal length $=2\left(\beta_{1} \beta_{2}, \beta_{1}^{-1} \beta_{2}\right.$ and $\beta_{2} \beta_{6}^{-1}$, which occur in $r$ and in $r^{-1}$ ), therefore it satisfies the small cancellation conditions $C^{\prime}(1 / 8)$ and $C(12)$.
Open question Do these properties of $\pi_{1}\left(S^{3} \backslash S_{W}\right)$ lead in general to some
small cancellation conditions, so that one can always solve the problem of deciding whether $i_{*}^{+}$is surjective?

## 5. - Appendix: The Alexander polynomial.

The Alexander polynomial is a well-known invariant for knots and links. Definitions and results about it can be found e.g. in [1, 18]. We are interested in the reduced Alexander polynomial for a closed braid, that is the polynomial in which the same variable $t$ has been associated to all components of the link.

Our algorithm, implemented in GAP, is adapted to links presented as closed braids written in band generators. Since the classical Artin generators form a subset of the band generators, it can be used also for braids written in the classical way.

The algorithm is based on the method of Seifert matrices:
Definition 14. - If S is a connected Seifert surface for a link $L$, define a Seifert matrix $M$ as follows: let $a_{1}, a_{2}, \ldots a_{N}$ be free abelian generators of $H_{1}(S)$ (here $N=1-\chi(S)$ ); then the elements of $M$ are

$$
M_{i j}=l k\left(a_{i}, i_{*}^{+}\left(a_{j}\right)\right),
$$

where lk denotes the linking number, $i^{+}$is the push-off map, and the $i_{*}^{+}\left(a_{j}\right)$ 's are elements of $H_{1}\left(S^{3} \backslash S\right)$.

Proposition 15. - If $M$ is any Seifert matrix for a given link $L$, then the Alexander polynomial is given by:

$$
\begin{gathered}
\Delta_{L}(t)=\operatorname{det}\left(t M-M^{T}\right) \text { if } L \text { has one component; } \\
\Delta_{L}(t)=\frac{\operatorname{det}\left(t M-M^{T}\right)}{1-t} \text { if } L \text { has more than one component. }
\end{gathered}
$$

In order to find a Seifert matrix for a braid $\beta$ given by a word $W$, we need to find a connected Seifert surface for it. The fact that the associated discband surface $S_{W}$ is connected corresponds to the fact that the associated graph $\Gamma_{W}$ is connected. If it is not connected, it is easy to find a connected one, by addying the necessary edges. This can be done on $W$ by conjugating it by the letters corresponding to the added edges. Our GAP procedure Connect$\operatorname{Braid}(W, n)$ gives a new word $W^{\prime}$ with connected $S_{W^{\prime}}$ in case $S_{W}$ was not connected.

We already know how to find generators of $\pi_{1}\left(S_{W}\right)$ in terms of edges $e_{k}$ of $\Gamma_{W}$ (see Section 3). The same loops $\alpha_{j}$ 's (or their abelianisations $a_{j}$ 's) can be used as generators of $H_{1}\left(S_{W}\right)$. We also already know how to find their images


Fig. 9. - How to compute $l k\left(a_{i}, i_{*}^{+}\left(a_{j}\right)\right)$.
$i_{*}^{+}\left(a_{j}\right)$ in $H_{1}\left(S^{3} \backslash S\right)$, for instance by abelianising the images $i_{*}^{+}\left(\alpha_{j}\right)$ found in Section 3.

It is now very easy to read the elements $M_{i j}$ (look at Figure 9):

$$
l k\left(a_{i}, i^{+}\left(a_{j}\right)\right)=\sum_{k=1}^{l}-\left(\text { power of } e_{k} \text { in } a_{i}\right)\left(\text { power of } b_{k} \text { in } i_{*}^{+}\left(a_{j}\right)\right)
$$

This is computed by our GAP procedure SeifertMatrix ( $W, n$ ).
The number of components of a closed braid can be easily read from $W$ :
$|\widehat{\beta}|=$ number of disjoint cycles in the induced permutation $\varrho(W) \in S_{n}$
(where $S_{n}$ is the symmetric group on $n$ elements). This is an invariant of the closed braid. Our GAP procedure IsnCycle $(p, n)$ finds «true» if $p$ is a permutation of just one cycle on $n$ points.

All the computations are performed in one step by our GAP procedure Alexander ( $W, n$ ).

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